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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
09/695,756	10/24/2000	Thomas W. Voshell	500080.02	2589

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EXAMINER

LAMARRE, GUY J

ART UNIT	PAPER NUMBER
2133	17

DATE MAILED: 04/24/2003

Please find below and/or attached an Office communication concerning this application or proceeding.

<b>Office Action Summary</b>	Application No.	Applicant(s)
	09/695,756	VOSHEL
	Examiner	Art Unit
	Guy J. Lamarre, P.E.	2133

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

**Period for Reply**

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If the period for reply specified above is less than thirty (30) days, a reply within the statutory minimum of thirty (30) days will be considered timely.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133).
- Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

**Status**

1) Responsive to communication(s) filed on 3 Feb. 2003.

2a) This action is **FINAL**.                    2b) This action is non-final.

3) Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

**Disposition of Claims**

4) Claim(s) 41-67 is/are pending in the application.

4a) Of the above claim(s) \_\_\_\_\_ is/are withdrawn from consideration.

5) Claim(s) \_\_\_\_\_ is/are allowed.

6) Claim(s) 41-67 is/are rejected.

7) Claim(s) \_\_\_\_\_ is/are objected to.

8) Claim(s) \_\_\_\_\_ are subject to restriction and/or election requirement.

**Application Papers**

9) The specification is objected to by the Examiner.

10) The drawing(s) filed on \_\_\_\_\_ is/are: a) accepted or b) objected to by the Examiner.

Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).

11) The proposed drawing correction filed on \_\_\_\_\_ is: a) approved b) disapproved by the Examiner.

If approved, corrected drawings are required in reply to this Office action.

12) The oath or declaration is objected to by the Examiner.

**Priority under 35 U.S.C. §§ 119 and 120**

13) Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).

a) All b) Some \* c) None of:

1. Certified copies of the priority documents have been received.

2. Certified copies of the priority documents have been received in Application No. \_\_\_\_\_.

3. Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

\* See the attached detailed Office action for a list of the certified copies not received.

14) Acknowledgment is made of a claim for domestic priority under 35 U.S.C. § 119(e) (to a provisional application).

a) The translation of the foreign language provisional application has been received.

15) Acknowledgment is made of a claim for domestic priority under 35 U.S.C. §§ 120 and/or 121.

**Attachment(s)**

1) Notice of References Cited (PTO-892)

2) Notice of Draftsperson's Patent Drawing Review (PTO-948)

3) Information Disclosure Statement(s) (PTO-1449) Paper No(s) \_\_\_\_\_.

4) Interview Summary (PTO-413) Paper No(s) \_\_\_\_\_.

5) Notice of Informal Patent Application (PTO-152)

6) Other: \_\_\_\_\_

## DETAILED ACTION

### Continued Examination Under 37 CFR 1.114

1. A request for continued examination (paper # 14) of 3 Feb. 2003, under 37 CFR 1.114, including the fee set forth in 37 CFR 1.17(e), was filed in this application after final rejection. Since this application is eligible for continued examination under 37 CFR 1.114, and the fee set forth in 37 CFR 1.17(e) has been timely paid, the finality of the previous Office action has been withdrawn pursuant to 37 CFR 1.114. Applicant's pre-amendment (paper # 16), concurrently filed, along with the request for 1-month extension of time(paper # 15), has been entered. This office action is in response to Applicants' pre-amendment.

2. **Claims 41, 45-47 and 59-62** are amended. **Claims 41-67** remain pending.

2.1 The prior art rejections of record are withdrawn in response to Applicants' amendment, filed on 31 Dec. 2002.

### Claim Rejections - 35 USC ' 103

3. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

3.0 **Claims 41-43, 45-48 and 54-64** are rejected under 35 U.S.C. 103(a) as being unpatentable over **Tsukakoshi** (US Patent No. 5,363,382; Nov. 13, 1991) in view of **Hancu et al.** "A concurrent test architecture for massively parallel computers and its error detection capability; IEEE Transactions on Parallel and Distributed Systems, Page(s): 1169 –1184, Vol. 5, Issue: 11; Nov, 1994."

**As per Claims 41, 45, 54-55, 59, 64, Tsukakoshi** substantially discloses the procedure for the claimed method of accessing a memory via test instructions or commands, means to initiate test mode, comprising: comparing memory addresses for a fault memory analysis (access/test) by performing substitute address allocation (mapping or remapping) and provides compression

means in col. 2 lines 35-50. **Not specifically described** in detail in **Tsukakoshi** is the step whereby compressing is effected on addresses as a means for reducing failure storage hardware.

**However** such memory address compressing approach is well known in memory testing. For example, **Hancu et al.**, in an analogous art, discloses algorithms “A concurrent test architecture” wherein such techniques are described including error detection routine based on comparison of routing or address signatures in compressed format. {See **Hancu et al.**, Id., Figs. 4-14, Tables 1-4 and associated description on pages 1170-1183, including intermediate queuing or buffering or memory means on page 1172 col. 2 para. 2. Data transfer performing equivalent functionality as data storage.} **Therefore**, it would have been obvious to a person having ordinary skill in the art at the time the invention was made to modify the procedure of **Tsukakoshi** by including therein address compression means as taught by **Hancu et al.**, because such modification would provide the procedure disclosed in **Tsukakoshi** with a technique whereby “*hardware is optimized so as to limit the size of the processing units for maximum integration and minimum cost.*” {See **Hancu et al.**, Id., page 1171 col. 2 para. 6 et seq.}

**As per Claim 42, 46-47 and 56, Hancu et al.** discloses the procedure for decompressing addresses for comparison so as to ease localization of failure on page 1172 col. 2 paras. 3-5

**As per Claims 43 and 48, Hancu et al.** discloses the procedure for decompressing addresses for comparison so as to ease localization of failure on page 1172 col. 2 paras. 3-5.

**As per Claim 58, Hancu et al.** discloses the procedure for initiating test command at powering up the memory on page 1169 col. 2 paras. 1-2, wherein errors are detected on-line or concurrently with execution of parallel program.

**As per Claim 60, Hancu et al.** discloses the procedure for intermediate queuing or buffering or memory or storage means on page 1172 col. 2 para. 2.

**As per Claim 61, Hancu et al.** discloses the procedure for testing or detecting faults prior to forwarding data to intermediate location or accepting memory requests on page 1172 col. 2 para. 2, page 1176 col. 1 Section B.

**As per Claim 62, Hancu et al.** discloses the procedure for compressing addresses prior to storage or transfer, e.g., on page 1169 col. 2 penultimate para., or page 1172 col. 2 para. 2, viz., "*both the routing and control streams are signatured (compressed).*"

**As per Claims 63 and 57, Hancu et al.** discloses the procedure for compressing addresses wherein all compressed addresses have unique signature to thereby ease localization of failure on page 1172 col. 2 paras. 3-5 subsequent to decompressing compressed addresses for comparison.

**3.1 Claims 44, 49-53, 65-67** is/are rejected under 35 U.S.C. 103(a) as being unpatentable over **Tsukakoshi** (US Patent No. 5,363,382; Nov. 13, 1991) in view of **Hancu et al.** "A concurrent test architecture for massively parallel computers and its error detection capability; IEEE Transactions on Parallel and Distributed Systems, Page(s): 1169 –1184, Vol. 5, Issue: 11; Nov, 1994" in further view of **Meaden** (US Patent No. 4,642,793; March 19, 1984).

**As per Claims 44, 49-53, 65-67, Tsukakoshi and Hancu** substantially disclose the procedure for the claimed method of claim 42. **Not specifically described** in detail in **Tsukakoshi** or **Hancu** is the step whereby calculating a value from the memory address comprises dividing the value represented by the memory address by a prime number or use of hashing code or function.

**However** such memory address generation approach is well known in data compression. For example, **Meaden**, in an analogous art, discloses algorithms in "*Many-to-one mapping hash address generator*" wherein such techniques are described. {See **Meaden**, Id., Abstract.} **Therefore**, it would have been obvious to a person having ordinary skill in the art at the time the invention was made to modify the procedure of **Tsukakoshi** and **Hancu** by including therein address generation means via prime numbers or hashing code or function as taught by **Meaden**,

because such modification would provide the procedure disclosed in **Tsukakoshi** and **Hancu** with a technique whereby "*The indicator R is also applied to the address input of a random access memory 37 having four locations, each of which contains a prime number in the range 3-251. The contents of the addressed location of the memory 37 supply the hashing key K for the hash coding circuit 30.*" {See **Meaden**, col. 3 line 27 et seq.}

**3.1.1 Claims 44, 49-53, 65-67** is/are rejected under 35 U.S.C. 103(a) as being unpatentable over **Tsukakoshi** (US Patent No. 5,363,382; Nov. 13, 1991) in view of **Hancu et al.** "A concurrent test architecture for massively parallel computers and its error detection capability; IEEE Transactions on Parallel and Distributed Systems, Page(s): 1169 –1184, Vol. 5, Issue: 11; Nov, 1994" in further view of **Matsuda** (US Patent No. 5,659,737; August 1, 1995).

As per **Claims 44, 49-53, 65-67**, **Tsukakoshi** and **Hancu** substantially disclose the procedure for the claimed method of claim 42. Not specifically described in detail in **Tsukakoshi** or **Hancu** is the step whereby calculating a value from the memory address comprises dividing the value represented by the memory address by a prime number or use of hashing code or function.

However such memory address generation approach is well known in data compression. For example, **Matsuda**, in an analogous art, discloses algorithms in "*Methods and apparatus for data compression that preserves order by using failure greater than and failure less than tokens*" wherein such techniques are described. {See **Matsuda**, Id., Abstract.} Therefore, it would have been obvious to a person having ordinary skill in the art at the time the invention was made to modify the procedure of **Tsukakoshi** and **Hancu** by including therein address generation means via prime numbers or hashing code or function as taught by **Matsuda**, because such modification would provide the procedure disclosed in **Tsukakoshi** and **Hancu** with a technique whereby "*FIG. 2 conceptually illustrates a mapping of predictive substrings, PS, from the character string 105 into the hash table 110. As shown by the lines and arrows in FIG. 2, for each predictive substring (PS), a hash index, specified by a hash function, maps the (PS) into an*

*entry in the hash table 110. For the example shown in FIG. 2, the predictive substring (PS.sub.i) containing the characters "ABC" maps into a hash table entry for storage of the successive character "D." Similarly, for the additional substrings (PS.sub.i+1, PS.sub.i+2, and PS.sub.i+3), each predictive substring maps into the hash table 110. These examples assume that the natural language indications properly predict the successive character. The order of the characters in hash table 110 is merely exemplary, and the actual storage of characters is based on the hash function. Any hash based predictive function may be used in conjunction with the present invention to map the predictive substrings into the entries of hash table 110. In one embodiment, a remainder of division technique is used. The remainder of division technique is defined by the function  $h(key) = key \bmod M$  where the divisor M determines the effective size of the hash table and is a prime number. The key is defined as the predictive substring. The remainder of division hash function works well when the block size is relatively small, such as when the block size is 3." {See Matsuda, col. 2 line 56 et seq.}*

**3.1.2 Claims 44, 49-53, 65-67** are rejected under 35 U.S.C. 103(a) as being unpatentable over **Tsukakoshi** (US Patent No. 5,363,382; Nov. 13, 1991) in view of **Hancu et al.** "A concurrent test architecture for massively parallel computers and its error detection capability; IEEE Transactions on Parallel and Distributed Systems, Page(s): 1169 –1184, Vol. 5, Issue: 11; Nov, 1994" in further view of **NN8806199** (Improved Hash and Index Searching Techniques for Computers Using a Cache And/Or Virtual Memory; IBM Technical Disclosure Bulletin, June 1988, US; VOLUME NUMBER: 31, PAGE NUMBER: 199 – 202, hereinafter **IBM Tech**).

As per **Claims 44, 49-53, 65-67**, **Tsukakoshi** and **Hancu** substantially disclose the procedure for the claimed method of claim 42. Not specifically described in detail in **Tsukakoshi** or **Hancu** is the step whereby calculating a value from the memory address comprises dividing the value represented by the memory address by a prime number or use of hashing code or function.

However such memory address generation approach is well known in data compression. For example, **IBM Tech**, in an analogous art, discloses algorithms wherein such techniques are

described. {See **IBM Tech**, Id., Abstract.} Therefore, it would have been obvious to a person having ordinary skill in the art at the time the invention was made to modify the procedure of **Tsukakoshi** and **Hancu** by including therein address generation means via prime numbers or hashing code or function as taught by **IBM Tech**, because such modification would provide the procedure disclosed in **Tsukakoshi** and **Hancu** with a technique whereby "*many subsets of the chained blocks are arranged so that they are contained in the same memory segment and their relative order is maintained. The standard index table contains an entry for each block comprising the array. The modified indexing table contains an entry for each memory segment that contains a logically contiguous subset of the chained blocks. In addition, each table entry contains a count of the number of blocks mapped by the entry, that is, the number of blocks contained within the memory segment indicated by the entry. If many blocks fit within a single memory segment, each entry in the indexing table may indicate several blocks instead of just one. This table is more compact than a general index that pointed to each block and therefore will cause less page and cache faults, in addition to occupying less memory.* The modified indexing structure is shown in Fig. 3. The same set of blocks as used in Fig. 2 are modified into a smaller indexing table." {See **IBM Tech**, last 10 lines.}

### Conclusion

- 4.1 The prior art made of record and relied upon is considered to applicant's disclosure.
- 4.2 Any response to this action should be mailed to:

Commissioner of Patents and Trademarks, Washington, D.C. 20231

or faxed to:

(703) 746-7238, (for After-Final communications),

(703) 746-7239, (for formal communications intended for entry),

Art Unit: 2133

(703) 746-5463 (for informal or draft communications, please label  
"PROPOSED" or "DRAFT").

Hand-delivered responses should be brought to Crystal Park II, 2121 Crystal Drive,  
Arlington, VA, Fourth Floor (Receptionist).

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Guy J. Lamarre, P.E., whose telephone number is (703) 305-0755. The examiner can normally be reached on Monday to Friday from 9:30 AM to 6:00 PM.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Albert De Cady, can be reached on (703) 305-9595.

Any inquiry of a general nature or relating to the status of this application or proceeding should be directed to the Group receptionist whose telephone number is (703) 305-3900.

Guy J. Lamarre, P.E.

*Guy J. Lamarre*

Patent Examiner

17 April 2003

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